A Cooperative Channel Estimation Approach for Coordinated Multipoint Transmission Networks

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Abstract—This paper considers the problem of coordinated multi-point transmission from transmitter devices that acquire noisy and partial channel state information (CSI) estimates from local feedback links. As cooperative multi-antenna transmission often relies on the availability of global channel estimates, we propose a novel decentralized algorithm that produces MMSE-optimal channel estimates on the basis of combining local feedback and inter-transmitter communications. To this end, we assume the devices are equipped with rate-limited bi-directional communication links over which they exchange a finite number of CSI-related bits. We propose a low-complexity cooperative channel estimation algorithm which exploits the local communications near-optimally and is robust to arbitrary feedback noise statistics. The proposed method has application in future decentralized CoMP scenarios, where we show clear advantages over the conventional channel state information exchange mechanisms.

I. INTRODUCTION

Coordinated multi-point (CoMP) transmission methods such as network MIMO [1] [2] [3], coordinated beamforming and scheduling [4] [5] [6], and interference alignment [7] often require the acquisition of global channel state information (CSI) at all the devices that engage in the coordination [8]. In practice, in Frequency Division Duplex (FDD) scenarios, CSI acquisition at the transmitter exploits local feedback links that provide at best a noisy and partial estimate of the total (network-wide) CSI matrix at each base station. In the current LTE release for instance, only the CSI related to a subset of the users served by one base station are acquired by that base station. In turn the cooperating transmitters can rely on signaling interfaces (such as X2 [9] alike interface in LTE) that provide limited inter-device communication capabilities so as to enable global channel estimates, hence facilitate cooperation. Currently, such links are subject to a coarse design however, limited to communicating the CSI coefficients that are missing at one transmitter, with no regard the statistical quality of the local information already existing at that transmitter, and ignoring potential benefits of correlated channel estimates between two transmitters, i.e., the fact that two base station may both obtain estimates of the CSI related to a common user.

In this paper we recast this problem into a more general and systematic decentralized channel estimation problem with side information. We start off with transmitters having acquired an initial amount of CSI from local feedback links. It is essential to note that this initial local CSI can be of completely arbitrary nature in our framework, including scenarios such as limited feedback from an arbitrary subset of users, with arbitrary estimation error statistics. Our framework also includes so-called hierarchical CSI scenarios where some transmitters are called to play a role of master in a cluster of small-cells, collecting more information than surrounding slave transmitters.

Based on the given initial local CSI structure, we design a two-step decentralized channel estimation aiming at producing minimum mean squared error (MMSE) optimal estimates at all cooperating transmitters. The new approach involves the transmitters (1) exchanging with each other shaped quantization representations of their local channel data, under a rate constraint on the inter-transmitter communication. Secondly, (2) the transmitters construct a final channel estimate based on a suitable combination of their local CSI and the data received from surrounding devices. Interestingly, the steps (1) and (2) are intertwined but fortunately can be solved jointly. Finally, when the inter-device communication links are bidirectional and total amount of bits for all coordination links are constrained, an interesting question is how many bits of channel-related information will flow in each directions. In other words which transmitters shall teach more to other transmitters (resp. learn more from them) about the channel state.

It has to be noted that the problem of cooperative channel estimation is rooted in the information theoretic framework of network vector quantization [10]. Specifically, for two transmitters cooperation, it is the problem of lossy source coding with side information (a.k.a. Wyner-Ziv coding) [11]. When more transmitters are involved in the cooperation, it is a generalization of Wyner-Ziv coding with the decoder takes into account multiple instances of correlated and encoded information as well as the side information. The information theoretical bound and asymptotically bound-achieving quantizer design for Wyner-Ziv coding are well analyzed, for example, in [12], [13], [14], [15], [16] and [17].

In this paper, we propose a novel algorithm (referred to as coordination shaping). For the two transmitters cooperation scenario, the proposed algorithm outperforms the asymptotic optimal Wyner-Ziv coding algorithm [17] in the low rate region and the performance asymptotically achieves the Wyner-Ziv bound in the high rate region. The proposed algorithm works for multi-transmitter cooperation scenarios as well, which are generalizations of Wyner-Ziv coding, to jointly optimize all the involved quantizers. Finally, our optimization framework allows to find the optimal bits allocation to each coordination

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The channel independent estimation error $E^{(i)}$ satisfies $e^{(i)} = \text{vec}(E^{(i)}) \sim \mathcal{CN}(0, Q_i)$. The estimation errors for different TXs are assumed independent, i.e., $E\{e^{(i)} \cdot e^{(j)}^H\} = 0, \forall i \neq j$.

Note that this CSI feedback model is quite general and includes diverse scenarios from local feedback to global noisy feedback. Representing various CSI information structure can be done by varying the statistics (i.e., $Q_i$) at the various transmitters, since larger coefficients on some elements of $Q_i$ induces a reduced information for the corresponding channels at $\text{TX}_i$, inducing partial CSI. Note that some feedback designs may lead to each transmitter having some noisy information about all the user channels. For instance in the broadcast feedback scenario, each RX acquires its downlink CSI from all transmitters (using pilots), quantizes it, then broadcasts it on the uplink feedback channel. Since the link strength between RX and each TX varies for each TX-RX pair, each TX then receives a different noisy version of the same global CSI.

In Fig. 1, the cooperation information exchange between two transmitters $\text{TX}_i$ and $\text{TX}_k$ are illustrated. $\text{TX}_k$ send to $\text{TX}_i$ a suitably quantized version of his local CSI estimate $\hat{h}^{(k)}_i$ through the coordination link, denoted here by $z_{ki}$. Similar operation is performed by $\text{TX}_i$ to provide $\text{TX}_k$ $z_{ik}$. Based on the local CSI and exchanged CSI, both transmitters compute a final estimate.

The rest of the paper is arranged as follows. Section II introduces the system model and describes the problem. Section III provides the optimization of the reconstruction function and Section IV performs the optimization of the quantizer. Section V solves the coordination link bit allocation problem and Section VI performs the optimization of the quantizer. The construction function is defined as $Q_ki : \mathcal{C}^{NL \times MK} \mapsto \mathcal{C}_k, \ z_{ki} \in \mathcal{C}_k, \ |\mathcal{C}_k| = 2^{R_{ki}}$, where $\mathcal{C}_k$ is the codebook for the quantizer $Q_ki$. Finally $R_{ki}$ is the finite number of bits on the coordination link between $\text{TX}_k$ and $\text{TX}_i$. All cooperation information is exchanged in a single shot and simultaneously with other transmitters.

The final estimate at $\text{TX}_i$ is denoted by $\hat{h}^{(i)}_i$, a reconstruction function $g_i(.)$ is used at $\text{TX}_i$ which combines local CSI $\hat{h}^{(i)}_i$ and exchanged CSI $z_{ki}$ to get $\hat{h}^{(i)}_i$. The quantization operation associated to the link from $\text{TX}_k$ to $\text{TX}_i$ is defined as $Q_{ki} : \mathcal{C}^{NL \times MK} \mapsto \mathcal{C}_k, \ z_{ki} \in \mathcal{C}_k, \ |\mathcal{C}_k| = 2^{R_{ki}}$, where $\mathcal{C}_k$ is the codebook for the quantizer $Q_{ki}$. Finally $R_{ki}$ is the finite number of bits on the coordination link between $\text{TX}_k$ and $\text{TX}_i$. All cooperation information is exchanged in a single shot and simultaneously with other transmitters.

The goal of this paper is to find (i) the optimal reconstruction function $g_i(.)$ and (ii) the optimal fix rate quantizer $Q_{ki}$ such that the expectation of the per dimensional MSE for the
final channel estimate

\[ D^{(i)} = \frac{1}{n} \mathbb{E}[\| \mathbf{h} - \hat{\mathbf{h}}^{(i)} \|^2] \]  

(2)
is minimized. Later on in this paper, we also optimize the values of \( R_{ki} \) under a global backhaul constraint.

Regarding the quantization on coordination link, the optimal vector quantization (VQ) is depicted in Fig. 2. \( z_{ki} \) and quantization error \( e_{Q_{ki}} \) are shown to be uncorrelated, \( e_{Q_{ki}} \) and \( \hat{\mathbf{h}}^{(k)} \) are dependent [18]. The covariance matrices for \( \hat{\mathbf{h}}^{(k)} \), \( z_{ki} \) and \( e_{Q_{ki}} \) satisfy \( \mathbf{Q}_{k}^{(k)} = \mathbf{Q}_{z_{ki}} + \mathbf{Q}_{e_{Q_{ki}}} \). Since the input of the quantizer \( \hat{\mathbf{h}}^{(k)} \) is Gaussian, we obtain an upper bound of the impact of quantization by assuming that the quantization error \( e_{Q_{ki}} = \mathbf{h}^{(k)} - z_{ki} \) is also Gaussian distributed as \( e_{Q_{ki}} \sim \mathcal{CN}(0, \mathbf{Q}_{e_{Q_{ki}}}) \) [19]. Similar to [20] and based on the assumption for \( e_{Q_{ki}} \) as Gaussian, we can approximate the VQ procedure by a gain-plus-additive noise model (similar to the scalar quantizer case in [21]) shown in Fig. 3.

**Proposition 1.** Assume the quantization error \( e_{Q_{ki}} \sim \mathcal{CN}(0, \mathbf{Q}_{e_{Q_{ki}}}) \), the optimal vector quantization for \( \hat{\mathbf{h}}^{(k)} \sim \mathcal{CN}(0, \mathbf{Q}_{e_{Q_{ki}}}) \) can be approximated with the quantization result \( z_{ki} \) satisfying:

\[ z_{ki} = (\mathbf{I} + \mathbf{Q}_{k} - \mathbf{Q}_{e_{Q_{ki}}})(\mathbf{I} + \mathbf{Q}_{k} - e_{Q_{ki}})^{-1}(\mathbf{h} + e^{(k)}) + q_{ki} \]  

(3)

where \( q_{ki}, \mathbf{h} \) and \( e^{(k)} \) are uncorrelated random vectors, \( q_{ki} \sim \mathcal{CN}(0, (\mathbf{I} + \mathbf{Q}_{k} - \mathbf{Q}_{e_{Q_{ki}}})(\mathbf{I} + \mathbf{Q}_{k} - e_{Q_{ki}})^{-1}\mathbf{Q}_{e_{Q_{ki}}}) \).

**Proof:** This is just a trivial generalization of gain-plus-additive noise model for optimal scalar quantizer in [21].

![Fig. 2. Quantizer model for optimal vector quantization.](image)

![Fig. 3. Gain-plus-additive noise model to approximate the optimal vector quantizer.](image)

In the reminder of this work, both the design of reconstruction functions and optimal quantizers will be based on (3).

### III. RECONSTRUCTION FUNCTION DESIGN

This section addresses the aforementioned sub-problem of the optimal reconstruction function \( g_{i} \) design in general setting of multi-transmitter cooperation \( K \geq 2 \). For the reason of simplicity, we only consider the reconstruction function as a weighted linear combination of estimates at TXi.

Therefore, the final estimate at TXi is:

\[ \hat{\mathbf{h}}^{(i)} = \sum_{k \in A_i} \mathbf{W}_{ki} z_{ki} + \mathbf{W}_{ii} \hat{\mathbf{h}}^{(i)} \]  

(4)

where the set \( A_i \) contains the indices of TXs that are cooperating with TXi. Thus, according to (3), with equations (2) and (4) the optimal weight combining matrices can be derived.

**Proposition 2.** The optimal weight combining matrices \( \{ \mathbf{W}_{ki}, \mathbf{W}_{ii}, k \in A_i \} \) for the optimization problem:

\[ \mathbf{w}_{ki}, \mathbf{w}_{ii}, k \in A_i \quad \text{s.t.} \quad \hat{\mathbf{h}}^{(i)} = \sum_{k \in A_i} \mathbf{W}_{ki} z_{ki} + \mathbf{W}_{ii} \hat{\mathbf{h}}^{(i)} \]

(5)

can be obtained from:

\[ \begin{bmatrix} \mathbf{W}_{i_1} & \mathbf{W}_{i_2} & \ldots & \mathbf{W}_{i_{C_i}} \end{bmatrix} = \mathbf{Y}_i \mathbf{\Omega}_i^{-1}. \]  

(6)

The optimum per dimensional MSE attained is:

\[ D^{(i)} = \frac{1}{n} \text{tr}(\mathbf{I} + \mathbf{Q}_{i}^{-1}) + \sum_{k \in A_i} ((\mathbf{I} + \mathbf{Q}_{k})(\mathbf{I} + \mathbf{Q}_{k} - \mathbf{Q}_{e_{Q_{ki}}})^{-1}(\mathbf{I} + \mathbf{Q}_{k} - \mathbf{I})^{-1})^{-1} \]  

(6)

where \( \mathbf{Y}_i, \mathbf{\Omega}_i \) are given below, the set \( A_i \) has cardinal \( |A_i| = C_i \) and is denoted by \( A_i = \{ l_1, \ldots, l_{C_i} \} \).

\[ \mathbf{Y}_i = \begin{bmatrix} \mathbf{I} & \mathbf{A}^{H}_{i_1} & \cdots & \mathbf{A}^{H}_{i_{C_i}} \\ \mathbf{I} + \mathbf{Q}_{i} & \mathbf{A}^{H}_{i_1} & \cdots & \mathbf{A}^{H}_{i_{C_i}} \\ \mathbf{A}_{i_1} & \mathbf{P}_{i_1} & \cdots & \mathbf{A}_{i_1} \mathbf{A}^{H}_{i_{C_i}} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{A}_{i_{C_i}} & \mathbf{A}_{i_{C_i}} \mathbf{A}^{H}_{i_1} & \cdots & \mathbf{P}_{i_{C_i}} \end{bmatrix} \]

\[ \mathbf{\Omega}_i = \begin{bmatrix} \mathbf{I} + \mathbf{Q}_{i} & \mathbf{Q}_{i} & \cdots & \mathbf{Q}_{i} \\ \mathbf{Q}_{i} & \mathbf{I} + \mathbf{Q}_{i} - \mathbf{Q}_{e_{Q_{ki}}} & \cdots & \mathbf{Q}_{i} \\ \mathbf{Q}_{i} & \mathbf{Q}_{i} & \cdots & \mathbf{I} + \mathbf{Q}_{i} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{Q}_{i} & \mathbf{Q}_{i} & \cdots & \mathbf{I} + \mathbf{Q}_{i} \end{bmatrix} \]

**Proof:** Due to lack of space, we only provide a sketch for the proof, the details are provided in the full size version of this paper [22]. According to (3), substitute \( z_{ki} \) into (4) and then (2), set the partial derivatives \( \frac{\partial D^{(i)}}{\partial \mathbf{W}_{ii}} = 0, \frac{\partial D^{(i)}}{\partial \mathbf{W}_{ki}} = 0, k \in A_i \) and solve this equation system, we can get (5). Insert (5) into (4) and then (2), after a few simplification we can get (6).

### IV. QUANTIZER DESIGN

Based on the weight expressions above, we can now proceed with the task of optimizing the quantizer used for communicating CSI-related bits from one transmitter to another. The main goal is to optimize \( \mathbf{Q}_{e_{Q_{ki}}} \) such that \( D^{(i)} \) can be minimized.
First, it’s obvious that a conventional VQ will not provide the best $s_{ki}$, as it typically minimizes the distortion in a way that ignores the local CSI available at TX$i$.

As $Q_i$, $i = 1 \ldots K$ reflect the accuracy of local CSI and are assumed to be known to all TXs, the quantizer should allocate the quantization resource where most needed, i.e. in channel elements or directions that are known best by the TX$k$ and least known by TX$i$. Thus, we propose to use the weighted square error distortion: $d_{Q_k} = (x - y)^H B_{ki} (x - y)$ as the distortion measure of the quantizer $Q_{ki}$, where the positive definite weight matrix $B_{ki}$ is a variable to be optimized.

Importantly, we can calculate $Q_{ki}$ in the asymptotic case when the given coordination link rate $R_{ki}$ is sufficiently large, i.e. the large quantization rate regime.

**Proposition 3.** In the large quantization rate regime, for a quantizer with $S$ quantization level, the quantization error covariance for a $n$ dimensional complex random vector source $x \sim CN(0, \Gamma)$ using optimal weighted mean square error distortion based vector quantization with weight matrix $B$ is $Q_k = 2S^{-\frac{n}{2}} M(S_0) 2\pi (\frac{n + 1}{n})^{n+1} \det(\Phi) = B^{-1}$, where $M(S_0)$ is the minimal normalized second moment for all space-filling polytopes in $2n$ dimension, and $\Phi$ is defined as:

$$
\Phi = \frac{1}{2} \left[ \Re(B_i^\frac{1}{2} \Gamma B^\frac{1}{2}_i) \Im(B_i^\frac{1}{2} \Gamma B^\frac{1}{2}_i) \right],
$$

**Proof:** See Appendix.

we now exploit Proposition 3 in order to derive $Q_{ki}$:

$$
Q_{ki} = 2^{1-\frac{\delta ki}{2}} M(S_0) 2\pi \left(\frac{n + 1}{n}\right)^{n+1} \det(\Phi_{ki}) \frac{1}{n} B_{ki}^{-1}
$$

(7)

where

$$
\Phi_{ki} = \frac{1}{2} \left[ \Re(B_i^\frac{1}{2} (I + Q_k) B_{ki}^\frac{1}{2} ) \Im(B_i^\frac{1}{2} (I + Q_k) B_{ki}^\frac{1}{2} ) \right],
$$

According to (6), (7), we can now optimize the quantization taking side information into account by solely optimizing the value of $B_{ki}$. When $R_{ki}$ is sufficiently large, the optimal $B_{ki}$ can be solved as follows:

$$
\min_{B_{ki}, k \in A_i} \quad D^{(i)*}
\text{s.t.} \quad \det(B_{ki}) = 1, B_{ki} \succeq 0
$$

(8)

The constraints on $B_{ki}$ matrices ensure that the optimization is performed among quantizers with the same quantization distortion. Note that a rigorous convexity analysis for optimization problem (8) is left for [22], only an approximation is provided here. Since we are considering the coordination link rate $R_{ki}$ to be sufficiently large, use matrix inverse approximation on (6):

$$
D^{(i)*} = \frac{1}{n} \text{tr}(I + Q_k^{-1}) + \sum_{k \in A_i} \left( (I + Q_k)(I + Q_k - Q_{ki})(I + Q_k - I)^{-1} \right)^{-1}
$$

$$
\simeq \frac{1}{n} \text{tr}\left( \sum_{k \in A_i} (Q_k + Q_{ki})^{-1} + Q_k^{-1} + I \right)^{-1}
$$

$$
\simeq \frac{1}{n} \text{tr}\left( \sum_{k \in A_i} Q_k^{-1} + Q_k^{-1} + I - \sum_{k \in A_i} Q_k^{-1} Q_{ki} Q_k^{-1} \right)^{-1}
$$

Therefore, the objective function is approximated as a trace of the inverse of positive semi-definite matrix, which is a convex optimization problem.

Note that once the optimal weight matrix $B_{ki}$ is obtained, the codebook for optimal quantizer $Q_{ki}$ can be calculated based on Lloyd algorithm and a training set. The optimal weight matrices for estimation combine can be calculated according to (5).

V. COORDINATION LINK BIT ALLOCATION PROBLEM

An interesting consequence of the above analysis is the optimization of coordination where multiple transmitters can exchange simultaneously CSI-related information to each other under a global constraint on the coordination bits. The global optimization problem over all coordination links now becomes:

$$
\min_{B_{ki}, k \in A_i, i = 1, \ldots, K} \quad \frac{1}{K} \sum_{i=1}^{K} D^{(i)*}
\text{s.t.} \quad \det(B_{ki}) = 1, B_{ki} \succeq 0
\sum_{i=1}^{K} R_{ki} = R_{tot}, R_{ki} \in \mathbb{N}^+
D^{(i)*} \text{ defined in (6)}
$$

Due to the integer constraints on $R_{ki}$, this problem becomes non-convex optimization. However, conventional alternating algorithms can be applied to perform the optimization in a two-step iterative approach.

VI. NUMERICAL PERFORMANCE ANALYSIS

In this section, the per dimensional MSE for decentralized channel estimation is evaluated for different settings using Monte-Carlo simulations over $10^6$ channel realizations. In all simulations $K = 2$ and $L = 2$, $M = N = 1$. The channel $h \in \mathbb{C}^{4 \times 1} \sim CN(0, I)$ and the rates on coordination link from TX2 to TX1 and from TX1 to TX2 are denoted $R_{21}$ and $R_{12}$, respectively. The parameter $M(S_0)$ is set as 929/12960 which is related to the ES lattice since we are considering 8 real dimensions. In Fig. 4 and Fig. 5, the local CSI information structure is characterized by $Q_1 = \text{diag}(0.1, 0.1, 0.9, 0.9)$ and $Q_2 = (0.9, 0.9, 0.1, 0.1)$ which corresponds to an example where TX1 has more accurate CSI about RX1 and less accurate CSI about RX2, while TX2 has more accurate CSI about RX2 and less accurate CSI about RX1. The $\text{diag}(\cdot)$ operator represents a diagonal matrix with diagonal elements in the parenthesis.

Fig. 4 shows the per dimensional MSE for the final estimation at TX1. The shaped coordination curve applies the
proposed algorithm. The unshaped coordination implement the traditional optimal VQ and find $W_{21}, W_{11}$ accordingly using (5). From the figure we can conclude that the shaped coordination algorithm outperforms the unshaped coordination algorithm, which not surprisingly shows the benefit of taking the priori statistic information into account. The asymptotically optimal Wyner-Ziv coding curve refers to the Wyner-Ziv quantizer for noisy source in [17]. It reveals that our algorithm outperforms the asymptotically optimal Wyner-Ziv coding algorithm in low coordination rate region. In fact, in two TXs cooperation case, it is trivial to prove that the proposed algorithm converges asymptotically to the Wyner-Ziv bound $D_\infty$ [17] as well.

Fig. 5 exhibits the sum rate for a 2 TX cooperation system. The rate on the coordination links satisfies $R_{21} = R_{12}$. The zero forcing (ZF) precoder is constructed at each TX based on its final estimate of the network MIMO channel. We also provide the sum rate for the case when coordination links have infinite bandwidth. The figure shows that the proposed shaped coordination algorithm will improve the system sum rate beyond the unshaped coordination algorithm and asymptotically optimal Wyner-Ziv coding algorithm when a simple ZF precoder is implemented. As the rate on coordination link increases, the sum rate for all algorithms will converge to the infinite coordination rate case.

Fig. 6 considers the coordination link bit allocation problem for a 2 TX cooperation system. The total amount of bits for coordination link is $R_{12} + R_{21} = 30$ bits, the error covariance matrices $Q_1 = diag(0.1, 0.1, 0.9, 0.9)$, $Q_2 = diag(0.9, 0.9, 0.1, 0.1)$. The figure reveals that the cooperation information exchange is not necessarily symmetric, the optimal coordination link bit allocation strategy is to let TX2 share the cooperation information to TX1 through a $R_{21} = 12$ bits coordination link and vice versa through a $R_{12} = 18$ bits coordination link. It reveals that if one TX has a better local CSI, it is more encouraged to share his information through a higher rate coordination link.

VII. CONCLUSION

We study the decentralized cooperative channel estimation for coordinated multipoint transmission networks. We derive a low-complexity algorithm which exploits local communications near-optimally and is robust to arbitrary feedback noise statistics. We exhibit clear advantages over CSI acquisition and exploitation methods used in conventional CoMP systems.

VIII. APPENDIX

Proof for Proposition 3: In order to prove this proposition, we need the following result: for a random vector $x$, if an optimal $S$ level Euclidean distance distortion based quantizer applied on a random vector $y = B^\top x$ has the codebook $\{y_1, \ldots, y_S\}$ and associated partition $\{P_1, \ldots, P_S\}$, then the optimal $S$ level weighted square error distortion based quantizer applied on the random vector $x$ with weight matrix $B$ will have the codebook $\{B^{-\frac{1}{2}}y_1, \ldots, B^{-\frac{1}{2}}y_S\}$ and associated partition $\{B^{-\frac{1}{2}}[P_1], \ldots, B^{-\frac{1}{2}}[P_S]\}$, where the $B^{-\frac{1}{2}}[P_i]$ is defined as $B^{-\frac{1}{2}}[P_i] = \{x : 3y \in P_i \text{ st. } x = B^{-\frac{1}{2}}y\}$. This result can be easily obtained by change of variables in the integral.
expression of the average quantizer distortion. Thus,
\[
Q_x = \mathbb{E}\{(x - Q(x))(x - Q(x))^H\} \\
= B^{-1/2} \mathbb{E}\{(y - Q(y))(y - Q(y))^H\} B^{-1/2} \\
= B^{-1} Q_y B^{-1/2},
\]
where \( y = B^2 x \sim \mathcal{CN}(0, B^2 \Gamma B^2) \). Let \( t = [\Re(y)^T \Im(y)^T]^T \), then \( t \sim \mathcal{N}(0, \Phi) \) and
\[
\Phi = \frac{1}{2} \begin{bmatrix}
\Re(B^2 \Gamma B^2) & \Im(B^2 \Gamma B^2) \\
\Im(B^2 \Gamma B^2) & \Re(B^2 \Gamma B^2)
\end{bmatrix}.
\]
Furthermore, it is proven in [23] that
\[
Q_t = \mathbb{E}\{(t - Q(t))(t - Q(t))^T\} = D_t I_n,
\]
where the average distortion \( D_t = \frac{1}{n} \text{tr}(Q_t) \) is obtained from [24] for large \( S \) as
\[
D_t = S^{-1/2} M(S_0) \left( \int f_t(t) \frac{n}{\pi \sigma^2} dt \right) \frac{n + 2}{n} \frac{1}{2} + \text{det}(\Phi) \frac{1}{n} \\
= S^{-1/2} M(S_0) 2\pi \left( \frac{n + 2}{n} \right)^{1/2} \frac{1}{2} + \text{det}(\Phi) \frac{1}{n} \\
\]
where \( f_t(.) \) is the probability density function (p.d.f) of \( t \). Finally, from the expression of \( Q_t \) and by a real-to-complex compression, we get
\[
Q_y = 2S^{-1/2} M(S_0) 2\pi \left( \frac{n + 2}{n} \right)^{1/2} \frac{1}{2} + \text{det}(\Phi) \frac{1}{n} \mathbf{I}_n
\]
which leads to
\[
Q_x = 2S^{-1/2} M(S_0) 2\pi \left( \frac{n + 2}{n} \right)^{1/2} \frac{1}{2} + \text{det}(\Phi) \frac{1}{n} B^{-1}.
\]

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